

Towards Quantitative Logics in Rocq

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Quantitative Logics (QLs) are logics whose truth values go beyond 0 and 1. They were first studied by Łukasiewicz [33] and Gödel [16, 20, 26], and have recently gained importance due to their applicability to safer machine learning [14]. A well-known and long-established important instance of QLs are fuzzy logics [11, 32] dealing with inherently vague propositions like “ x is much larger than 10”, valued in the interval $[0, \infty]$. Much more recently, Bacci et al. [5–7] present propositional QLs, for example to reason about inequalities between rational functions. They are valued in structures based on the Lawvere quantale [23], with $\wedge$ and $\vee$ interpreted by max and min. Bacci and Møgelberg [8] introduce a higher-order QL interpreted in $[0, 1]$, newly allowing for unified reasoning on, but not limited to, program logics for probabilistic programming languages, bisimilarity on probabilistic processes, and properties of machine learning. Lastly, Capucci et al. [10] propose a QL parametric in $p \in (0, \infty]$ with quantitative proof calculus, called *p-hard quantitative linear logic* (*pQLL*).

Quantitative logics are becoming increasingly important, yet they lack a general formal theory. Therefore, we are interested in establishing a solid formal basis for QLs, envisaging a general and reusable library for these logics mechanised in Rocq [28]. As a key ingredient to many QLs are intervals of (extended) real numbers, our implementation should allow for extensive reasoning about these numbers at its core. We will be building on top of an exploration of selected traditional and recent propositional QLs in Rocq by Affeldt et al. [3, 4], also taking inspiration from a position paper by Marulanda-Giraldo et al. [25]. As a case study, we work towards mechanising *pQLL*.

In this abstract, we discuss

1. the ongoing effort to build a theory of nonnegative extended real numbers in Rocq, including connectives often used in QLs, relying on MathComp [24] and MathComp-analysis [1],
2. key features of *pQLL* and their mechanisation, alongside work towards a novel completeness result, and
3. improvements made and being made to MathComp(-analysis) as part of the process.

Our mechanisation can be accessed on <https://github.com/Janis-Bai/QLLib/releases/tag/Rocqshop2026>. The links in the document point to the corresponding Rocq code.

Extended real numbers in Rocq. Many QLs are based on intervals of real numbers such as $[0, \infty]$ used in *pQLL*. Besides standard operations such as **multiplication** $\otimes$, we require the notions of *comultiplication* $\otimes^*$ and *p-sums* $\oplus^p$. **Comultiplication** $a \otimes^* b := (a^{-1} \otimes b^{-1})^{-1}$ only differs from multiplication for $a = 0$ and $b = \infty$ (and vice versa), and *p-sums* are a soft variant of max and min, explained below.

For machine learning applications, it is desirable to have operations that are differentiable and componentwise strictly increasing (*smooth monotone*) [29, Property 4], also referred to as *soft operations*. The motivation behind this stems from learning algorithms such as gradient descent, which solve a minimisation problem to find optimal network parameters. For a soft operation, a small change in one of the parameters will result in a change in the overall value, improving the algorithm’s effectiveness. This is contrary to more traditional non-soft operations such as max and min, which are referred to as *hard operations*.

Hence, due to being soft, *p-sums* $a \oplus^p b := (a^p + b^p)^{1/p}$ and **harmonic p-sums** $a \oplus^{-p} b := (a^{-1} \oplus^p b^{-1})^{-1}$, originally applied to QLs in Yager logic [31], have recently gained importance and are used, for example, in *pQLL*. A key relation we are currently mechanising is that $\oplus^p$ converges to the binary maximum function as $p \rightarrow \infty$.

As of now, we mechanised all the operations above, along with a few fundamental facts such as **associativity of $\oplus^p$** .

Quantitative linear logic. In *pQLL*, there are soft connectives $\vee$ and $\wedge$ interpreted by $\oplus^p$ and $\oplus^{-p}$. This approach sacrifices idempotence (since $1 \oplus^p 1 = 2^{1/p} \neq 1$), which is related to a general result about logical connectives: connectives cannot be associative, idempotent, and smooth monotone at the same time [29, Proposition 1]. Also, the algebraic and proof-theoretic properties of QLs with soft connectives

are not well-studied, unlike those with hard connectives, which have an extensive theory since the 1990s, e.g. via residuated lattices [19, 30].

A key novelty of p QLL is its *quantitative proof calculus*, which is based on sequent calculus: deduction rules are annotated by operations on or constants in the extended reals. For example, the axiom rule and the left rule for additive soft disjunction are

$$\frac{1}{A \vdash A} \text{ AX} \qquad \frac{\Gamma \vdash A, \Delta \quad \oplus^p \quad \Gamma \vdash B, \Delta}{\Gamma \vdash A \vee B, \Delta} \vee_R.$$

Given a proof, one can compute a unique extended real number by traversing the proof tree, called the *validity* of the proof. For instance, the proof of $A \vdash A \vee A$ using $\vee_R$ and then AX twice has validity $1 \oplus^p 1 = 2^{1/p}$. This notion can be extended to sequents $\Gamma \vdash \Delta$, called *provability* and denoted $|\Gamma \vdash \Delta|$, defined as $\sup\{\text{validity}(P) \mid P \text{ is a proof of } \Gamma \vdash \Delta\}$, noting that $\sup \emptyset = 0$. The mechanisation of these definitions is standard and outlined below.

Mechanisation of p QLL’s syntax and deduction system. Formulas in p QLL are given by the following grammar, where V is a type of propositional variables:

$$A, B : \mathcal{F}_V ::= a \mid a^* \mid 1 \mid A \otimes B \mid A \otimes^* B \mid \perp \mid \top \mid A \vee B \mid A \wedge B \qquad a \in V$$

We mechanised sequents using an inductive type with type constructor $\vdash : \text{List}(\mathcal{F}_V) \rightarrow \text{List}(\mathcal{F}_V) \rightarrow \text{Type}$ and one value constructor per deduction rule, ignoring the annotations referring to a proof’s validity. To compute a proof’s validity, we implemented a *fixpoint of type* $\forall \Gamma, \Delta. \Gamma \vdash \Delta \rightarrow [0, \infty]$ recursing on the derivation of $\Gamma \vdash \Delta$. *Provability* $|\Gamma \vdash \Delta|$ is mechanised by forming the set S of all proofs of $\Gamma \vdash \Delta$, and then taking the supremum of $\{\text{validity}(p) \mid p \in S\}$.

Towards completeness. Semantics for p QLL is defined via *softales* [10, Definition 5.2], a preorder-like structure where the order relation is quantitative, i.e. $a \leq b$ is not a binary truth value, but an extended real number. Using this notion, Capucci et al. prove that p QLL is sound [10, Proposition 5.9].

We are currently working towards a novel completeness result for 1QLL without propositional variables, and currently conjecture the following lemma, which we believe follows by induction on A .

Lemma 1. *For $A \in \mathcal{F}_\emptyset$, we have $\llbracket A \rrbracket \leq |\vdash A|$, where $\llbracket \cdot \rrbracket : \mathcal{F}_\emptyset \rightarrow [0, \infty]$ is the interpretation of formulas.*

From this lemma, we believe that the completeness result follows readily once softales are mechanised. We expect the latter to be straightforward using the Hierarchy Builder [13].

Implementation details. In Rocq, operations such as $\otimes^*$, $\oplus^p$ operate on the type $\{\text{nonneg} \setminus \text{bar } \mathbb{R}\}$, morally a Σ -type of an extended real number x and a certificate that $0 \leq x$.

The key challenge in defining these operations is to verify that they preserve nonnegativity. While MathComp has automation to infer these and other preservation properties, known as interval inference [2], this automation was not fully implemented and tested yet for extended reals. For example, during our mechanisation we found that the $+$ notation for addition on extended reals was not recognised by the inference automation. Further, we proved the required inference lemmas for *inversion* and *powers*.

We defined p -sums $a \oplus^p b$ as p -integral $\|f\|_p := (\int_\Omega |f|^p d\mu)^{-p}$, where μ is the counting measure on $\Omega := \mathbb{N}$ and $f(0) = a$, $f(1) = b$, and $f(n+2) = 0$. This approach may seem cumbersome at first, but it has the advantage of allowing results about p -integrals to be directly applied to p -sums. Additionally, some properties of p -sums can be mechanised in greater generality.

The original definition of p QLL uses finite multisets to define sequents, while in our mechanisation of the calculus, we employ the well-established type of lists because of its ease of automation and speed. As a consequence, exchange rules were added in Rocq which are absent in the original definition of the logic. In the future, we may facilitate an implementation of finite multisets based on MathComp [12], inspired by Affeldt et al. [4], who point out that multisets may allow for easier proof automation in the absence of exchange rules.

Future work. Cut elimination, which is admitted by p QLL [10, Theorem 4.1], enables proof search and facilitates the consistency proof for p QLL. Motivated by this, we are currently mechanising cut elimination, having already completed the *core argument* for a *one-sided variant* of p QLL subject to a *purely algebraic lemma*. We conjecture the single-sided calculus to be equivalent to p QLL. The cut elimination proof extends Laurent’s mechanised proof [22] for linear logic, following their high-level structure as well as the use of a one-sided calculus, and adding reasoning about proof validities as required by p QLL.

Adding quantification to p QLL is an ongoing effort [9, 10] motivated by its applications to machine learning [17], where specifications are often first-order. We are interested in mechanising this, which would require dealing with object-level binders. The Rocq Libraries for First-Order Logic [21] and Undecidability [18] might provide a good framework as their type of formulas is quantified over arbitrary logical atoms, connectives, and quantifiers and as these libraries provide extensive support for binders using de Bruijn syntax [15], inspired by Autosubst 2 [27].

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