

QuantumLib:

A Library for Quantum Computing in Coq





Collaborators





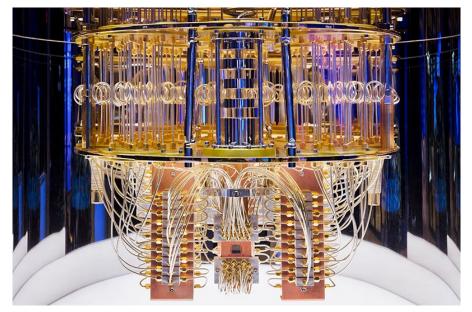


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Additional list of contributors can be found at: <u>https://github.com/inQWIRE/QuantumLib</u>

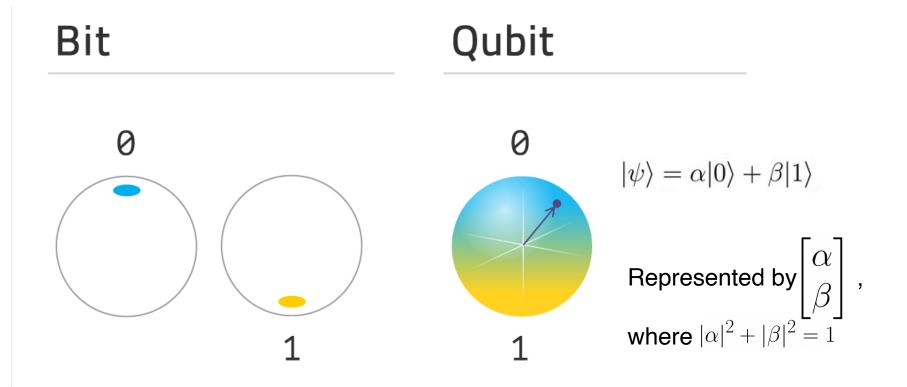
Why Verify Quantum Computing?

- Can be faster than classical systems
 - Quantum simulation
 - Shor's algorithm: encryption
 - Grover's search algorithm
- Quantum Advantage in practice
 - Google's random circuit sampling
 - Boson sampling (USTC, Xanadu)
- Quantum computing is hard!
 - Conceptually hard
 - Very error-prone



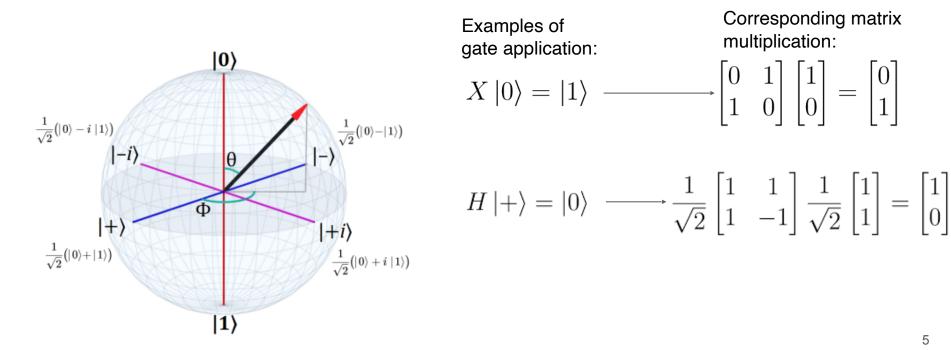
IBM's 127 qubit quantum computer

Quantum Bits: Qubits



Quantum Gates: Unitary Operators

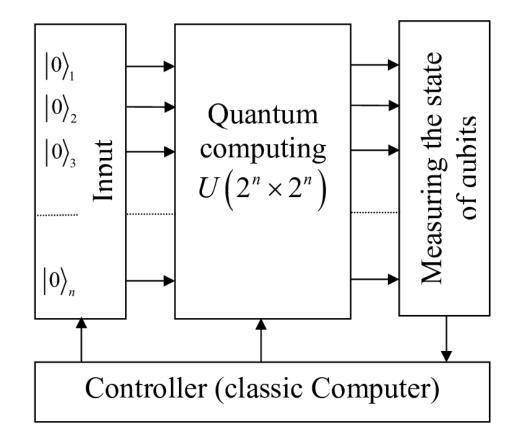
- Gates act on qubits to change their state
 - Eg: X, Y, Z, H, S, T 0



Quantum Circuits

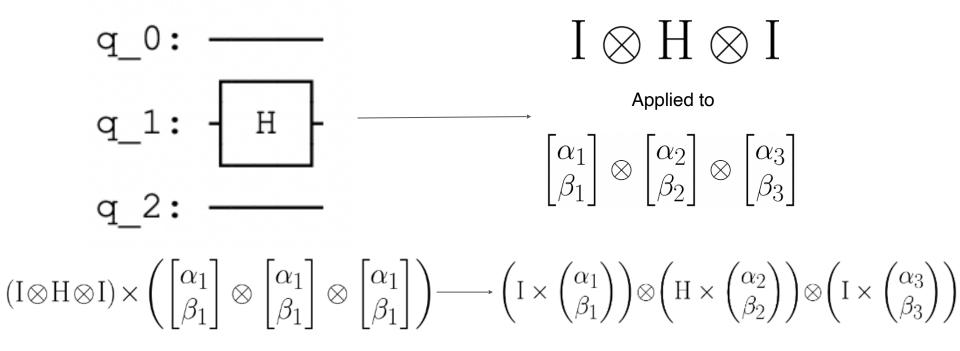
• Circuits represented by kronecker product:

 $\frac{\alpha_n}{\beta}$ α_1 $\otimes ... \otimes$



Quantum Programs

• Applying gates corresponds to multiplication by padded matrix:



Why QuantumLib?

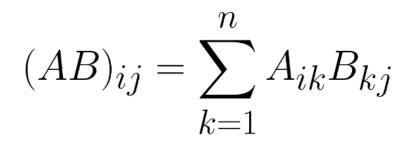
- Provide a backbone for quantum computing projects in Coq
- Tailored specifically towards quantum computing
 - QuantumLib is more efficient and comprehensive than other more general libraries
 - Can act as an extension of MathComp or Ccorn
- Consists of both low level and high level components

To do this, we rigorize the notions of Hilbert spaces in Coq: \mathbb{C}^{2^n}

Underlying Field Structure: Complex Numbers

- Coquelicot's complex numbers:
 - Added lemmas involving Euler's identity
- Polar coordinates
 - Eg: e^{ix}
- Summation notation
- Computable for our purposes
 - Sufficient rewrite lemmas
 - \circ Proof that \mathbb{C} is a field
- Polynomials over C and proof of completeness
 - Used for facts about determinants and eigenvectors

Definition C := (R * R)%type.



Matrices Over C

• Matrices defined as follows:

```
Definition Matrix (m n : nat) := nat -> nat -> C.
```

• Matrices must be well-formed:

Definition WF_Matrix {m n: nat} (A : Matrix m n) : Prop :=
 forall x y, x >= m \/ y >= n -> A x y = 0.

• Examples:

```
Definition \sigma x: Matrix 2 2 :=

fun x y => match x, y with

\begin{vmatrix} 0, 1 => 1 \\ 1, 0 => 1 \\ \\ \end{vmatrix}, \_ => 0

\sigma_x = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}

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Some Example Operations

```
Definition Mplus {m n : nat} (A B : Matrix m n) : Matrix m n :=
  fun x y => (A x y + B x y).
```

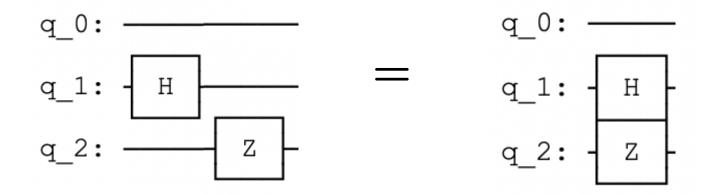
Definition Mmult {m n o : nat} (A : Matrix m n) (B : Matrix n o) : Matrix m o :=
fun x z => big_sum (fun y => A x y * B y z) n.

Applying circuits in series

```
Definition kron {m n o p : nat} (A : Matrix m n) (B : Matrix o p) :
    Matrix (m*o) (n*p) :=
    fun x y => (A (x / o) (y / p)) * (B (x mod o) (y mod p)).
```

Applying circuits in parallel

Compatibility between Mmult and kron



Applying circuits in series

Applying circuits in parallel

Lemma program_equivalence : $(I \otimes H \otimes I) \times (I \otimes I \otimes Z) = (I \otimes H \otimes Z)$.

Why These Design Choices?

- Phantom types help with proofs
 - Useful since kron changes size of matrix

```
Lemma kron_assoc : forall {m n p q r s : nat}
(A : Matrix m n) (B : Matrix p q) (C : Matrix r s),
(A \otimes B) \otimes C = A \otimes (B \otimes C).
```

Matrix ((m*p)*r) ((n*q)*s) Matrix (m*(p*r)) (n*(q*s))

Definition pad {n} (start dim : nat) (A : Square (2^n)) : Square (2^dim) :=
 if start + n <=? dim then I (2^start) @ A @ I (2^(dim - (start + n))) else Zero.</pre>

- Not clear that $(2^{\text{start}} * 2^n * 2^{\dim (\text{start} + n)}) = 2^{\dim n}$
 - Relies on guard

Other components

- Comprehensive linear algebra
 - Linear independence, diagonalizability, determinant
 - Eigenvectors
- Quantum components
 - Low level: basic states and gates
 - High level: padded gates, pure/mixed states
- Measurement and probability theory
- Translation to/from different representations of quantum objects
 - Vector states
 - Permutations

```
Fixpoint f_to_vec (n : nat) (f : nat -> bool) : Vector (2^n) :=
  match n with
  | 0 => I 1
  | S n' => (f_to_vec n' f) @ | f n' )
  end.
```

```
Lemma H0_spec : hadamard \times |0\rangle = |+\rangle.
Proof. lma'. Qed.
```

```
Lemma H1_spec : hadamard × |1\rangle = |-\rangle.
Proof. lma'. Qed.
```

Definition pad_u (dim n : nat) (u : Square 2) :
 Square (2^dim) := @pad 1 n dim u.

Matrix Tactics

- Well-formedness tactics
 - \circ show_wf and wf_db
- Ima
- gridify
 - Good for symbolic proofs
- solve_matrix
 - Good for numerical computations
- restore_dims

lma:

- Breaks down equality statement into cells
- Eg:

Lemma MmultYY : $\sigma y \times \sigma y = I 2$. Proof. lma. Qed.

gridify:

```
Lemma pad_mult : forall n dim start (A B : Square (2^n)),
pad start dim A × pad start dim B = pad start dim (A × B).
Proof.
    intros.         (I (2 ^ start) ⊗ A ⊗ I (2 ^ d)) × (I (2 ^ start) ⊗ B ⊗ I (2 ^ d))
    unfold pad.         =
    gridify.         I (2 ^ start) ⊗ (A × B) ⊗ I (2 ^ d)
    reflexivity.
Qed.
```

```
Lemma pad_A_B_commutes : forall dim m n A B,
m <> n ->
WF_Matrix A ->
WF_Matrix B ->
pad_u dim m A × pad_u dim n B = pad_u dim n B × pad_u dim m A.
Proof.
intros.
unfold pad_u, pad.
gridify; trivial.
Qed.
```

QuantumLib in Practice

- QWIRE
 - Quantum language built in QWIRE
 - Dependent types
- SQIR/VOQC
 - Uses translations between matrices
 - Proof of Shor's algorithm, Grover's algorithm, teleportation
- VyZX
 - ZX calculus: graphical representation of quantum computing
- Verification of Gottesman logic

Why QuantumLib Instead of Other Libraries?

- Works well with Kronecker product
- Tailored specifically to quantum computing
 - Contains additional components that other libraries lack
- Large math foundation
- Many rewrite lemmas/tactics
- Very compatible with other libraries (coming soon)

Challenges

- Slow when matrices get large
 - 3 qubits much worse than 2
 - 4 qubits very slow
 - Multiplication: $O(n^3)$ when n is matrix size so $O(2^{3n})$ when n is number of qubits
- Showing dimensions align
 - kron mixed product rule can be hard to apply
- Not actually computable
 - Leads to more slowness

Things for the Future

- Make tactics even better
 - Improve gridify by having different versions
- Add more linear algebra proofs
 - More matrix properties
 - Proof of Gottesman-Knill theorem
- Add more lemmas to Pad.v
 - More commutation lemmas
- Version with Ccorn for more computable matrices

More Things for the Future

- Group, ring, field typeclasses
 - field and ring exist in Coq, why not other structures?
 - Reification-style tactics
 - Generalize matrices to be over any field
 - Would make integration with Ccorn very easy
 - Finite fields

Summary

- Quantum computing occurs over Hilbert spaces
 - Strong interplay between matrix multiplication and kronecker product
- Matrix definition: Definition Matrix (m n : nat) := nat -> nat -> C.
 - Matrices must also be well formed
- Defined many quantum computing notions
- Matrix tactics
- Project found at https://github.com/inQWIRE/QuantumLib
- To install:
 - opam install coq-quantumlib